

This homework is due in class on Friday, November 16.

1. Give an English language description of the language generated by each of the following regular expressions over the alphabet $\{a, b, c\}$.

a) $(a + b)^*c(a + b)^*$

b) $(a + b + c)^*b(a + b + c)^*c(a + b + c)^*$

c) $(a + b + c)b(a + b + c)^*c(a + b + c)$

2. Give an argument to show that if L is a regular language, then L^R is also a regular language. (Hint: Start with a DFA that accepts L , and think about how you can produce a machine (either a DFA *or* and NFA) that accepts L^R .) You do not have to give a completely rigorous proof, but you should give a convincing argument.

3. Consider the following alphabet, in which each symbol is a stack of three 0's and 1's:

$$\Sigma = \left\{ \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix}, \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix}, \begin{bmatrix} 0 \\ 1 \\ 1 \end{bmatrix}, \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 1 \\ 0 \\ 1 \end{bmatrix}, \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix}, \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix} \right\}$$

The addition of two binary numbers can be represented as a string over this Σ , with each symbol in the string representing one of the columns in the addition. For example,

$$\begin{array}{r} 00100101 \\ 01001011 \\ \hline 01110000 \end{array} \text{ is represented as } \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \\ 1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \\ 1 \end{bmatrix} \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix}$$

A string over Σ , interpreted as an addition of binary numbers can be either a *correct* addition or an *incorrect* addition. Let L be the language over Σ that consists of all strings that represent correct additions. Show that this language is regular by constructing a DFA that accepts the language. If you prefer, you can show that L^R (and therefor also L) is regular by constructing a DFA that accepts L^R — it is actually easier to think about the problem when the DFA reads the string from right-to-left. This exercise shows that the problem of checking whether a binary addition is correct is one that can be solved by a DFA.

4. Use the pumping lemma to prove that the language $L_1 = \{a^n b^m c^k \mid k = n + m\}$ is not regular.
5. Use the pumping lemma to prove that the language $L_2 = \{w \in \{a, b\}^* \mid w^R = w\}$ is not regular.
6. Use the pumping lemma to prove that the language $L_3 = \{a^{2^n} \mid n \in \mathbb{N}\}$ is not regular.