The second test for this course will be given in class on Wednesday, October 20. It covers everything that we have done from Chapter 2, Sections 1, 2, 3, 4, and 6; it also covers Section 3.1. Note that we skipped some of the material in the listed sections in Chapter 2. As usual, the test will include some "essay-type" questions that ask you to define something, or discuss something, or explain something, and so on. Other than that, you can expect most of the questions to be similar to problems that have been given on the homework. This might include some proofs. There might be some questions related to the programming assignment and the various ways of representing sets in Java.

Here are some terms and ideas that you should be familiar with for the test:

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sets
set notations: \{a, b, c\}, \{1, 2, 3, \dots\}, \{x \mid P(x)\}, \{x \in A \mid P(x)\}
the empty set, \emptyset or \{\}
equality of sets: A = B iff they contain the same elements
element of a set: a \in A
sets can contain other sets as elements
subset: A \subseteq B
A = B if and only if both A \subseteq B and B \subseteq A
proper subset
union, intersection, and set difference: A \cup B, A \cap B, A \setminus B
definition of set operations in terms of logical operators
disjoint sets (A \cap B = \emptyset)
power set of a set: \mathcal{P}(A)
universal set
complement of a set (in a universal set): \overline{A}
DeMorgan's Laws for sets
Russell's Paradox
there is no "set of all sets"
bitwise operations in Java: &, |, ~
using an n-bit integer to represent subsets of \{0, 1, 2, \dots, n-1\}
&, |, and ~ as set operations (intersection, union, complement)
the shift operators <<, >>, and >>>
using "1 << n" to represent the singleton set \{n\}
representing a set of integers as a boolean array in Java
Java's TreeSet<Integer> class; operations set.add(n), set.remove(n), set.contains(n)
ordered pair: (a, b)
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equality of ordered pairs; (a, b) = (c, d) iff a = c and b = d
cross product of sets: A \times B
one-to-one correspondence
cardinality of a finite set: |A|
cardinality rules for finite sets: |A \times B| = |A| \cdot |B|, |\mathcal{P}(A)| = 2^{|A|}, |A \cup B| = |A| + |B| - |A \cap B|
finite set (in one-to one correspondence with one of the sets N_0, N_1, N_2, \dots)
infinite set (not finite)
countably infinite set (in one-to-one correspondence with \mathbb{N})
a set is countably infinite iff its elements can be placed into an infinite list
countable set (finite or countably infinite)
uncountable set (that is, uncountably infinite)
Cantor's dialgonalization proof that the set of real numbers is uncountable
examples of countably infinite and uncountably infinite sets
the union of two countably infinite sets is countably infinite
the cross product of two countably infinite sets is countably infinite
if X is uncountable and A is a countable subset, then X \setminus A is uncountable
for any set A, there is no one-to-one correspondence between A and \mathcal{P}(A)
proof of the above fact (given f: A \to \mathcal{P}(A), consider X = \{a \in A \mid a \notin f(a)\})
the power set of a countably infinite set is uncountable
function f: A \to B
a function f: A \to B as a subset of A \times B
alphabet (finite, non-empty set of "symbols")
string over an alphabet \Sigma
length of a string, |x|
empty string, \varepsilon
concatenation of strings, xy or x \cdot y
reverse of a string, x^R
x^n, for a string x and a natural number n
n_{\sigma}(x), the number of occurrences of a symbol \sigma in a string x
the set of all possible strings over \Sigma, denoted \Sigma^*
language over an alphabet \Sigma (a subset of \Sigma^*)
a language over \Sigma is an element of \mathcal{P}(\Sigma^*)
the set of strings over \Sigma is countable; the set of languages over \Sigma is uncountable
operations on languages
union, intersection, set difference, and complement applied to languages
concatenation of two languages: LM
L^n, for a language L and a natural number n
Kleene star of a language: L^*
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