This homework covers sections 3.7, 4.1, and 4.2, with a little left over from section 3.6. It is due in class Wednesday, April 10. Hand in a hardcopy of your solutions.

While you may discuss problems with other students, you should always make the first attempt on a problem yourself and you must write up your own solutions in your own words. You may not collaboratively write solutions or copy a solution that one person in the group writes up.

1. Use the construction outlined in the proof of Theorem 3.5 to construct a machine that accepts the intersection of the languages accepted by the two NFAs shown. (Hint: remember that the construction outlined in the proof of Theorem 3.5 is for two *DFAs*.)



2. Use the Pumping Lemma (Theorem 3.6) to prove that the following languages are not regular.

Remember the pattern for the proof discussed in class, which can be written:

Suppose L is regular. Then, by the Pumping Lemma, there is an integer N such that for any string  $w \in L$  with  $|w| \geq K$ , w = xyz where  $|xy| \leq N$ ,  $|y| \geq 1$ , and  $xy^n z \in L$  for all natural numbers n. But let w =\_\_\_\_\_, which is in L with  $|w| \geq N$ . Write w as xyz. Because  $|xy| \leq N$ , y must be of the form \_\_\_\_\_. But then  $xy^n z \notin L$  for n =\_\_\_\_\_ because \_\_\_\_\_.

Fill in the blanks.

- (a)  $L = \{ a^n b^n c^n \mid n \in \mathbb{N} \}$
- (b)  $L = \{ w \in \{a, b\}^* \mid n_a(w) < n_b(w) \}$ , where  $n_\sigma(w)$  means the number of  $\sigma$ s in w

3. Consider the context-free grammar shown on the right.

(a)	Write a derivation for	the string $aabbc$ using this grammar.	1.	$S \longrightarrow TR$
			this grammar.	$T \longrightarrow aTb$

- (b) Write a derivation for the string *abcccdd* using this grammar.  $T \longrightarrow \epsilon$
- (c) Find the language generated by this grammar. Briefly justify  $\begin{array}{c} R \longrightarrow cRd \\ R \longrightarrow c \end{array}$
- 4. For each of the following languages, give a context-free grammar that generates the language. Also explain in words why your grammar works.

Hint for part (d): keep in mind how to generate equal numbers of things, and think about which pairs of characters can match up with each other.

- a)  $\{a^n b^m \mid n \neq m\}$  b)  $\{a^n b^m c^k \mid m > n + k\}$
- c)  $\{a^{n}b^{m}c^{k}d^{l} \mid m = k \text{ and } n = l\}$  d)  $\{a^{n}b^{m}c^{k}d^{l} \mid n + m = k + l\}$
- 5. Given the following BNF grammar for a (very small) subset of Java syntax, write down six things generated by this grammar. Your examples should demonstrate all of the possibilities represented in the rules make sure to use each of the different rules and variations in at least one derivation.

$$\begin{array}{l} \langle name \rangle ::= \langle object\_ref \rangle \ ["." \ \langle identifier \rangle \ ] \\ \langle object\_ref \rangle ::= \langle identifier \rangle \ | \ \langle method\_call \rangle \\ \langle method\_call \rangle ::= \langle identifier \rangle \ "(" \ \langle name \rangle \ ["," \ \langle name \rangle \ ] \dots \ ")" \\ \langle identifier \rangle ::= "a" \ | "b" \ | "c" \ | "x" \ | "y" \ | "z" \end{array}$$

Write a BNF grammar for numeric literals in Java. Numeric literals include integers (e.g. 0, 1035, -47, +2) and floating point numbers (e.g. 17.3, 0.73, .50, 23.1e67, -1.34E-12, +0.2, 100E+100). Note that the only numbers that can start with 0 are the integer 0 and a floating point number between -1 and 1.