	Let $n$ be an arbitrary integer.	
Proof	1. <i>n</i> is even	premise
for all integers $n$ , if $n$ is even, then	<ol> <li>if n is even, then n = 2k for some integer k</li> <li>n = 2k for some integer k</li> <li>if n = 2k for some integer k,</li> </ol>	definition of even from 1, 2 (modus ponens)
n² is even	then $n^2 = 4k^2$ for that integer k 5. $n^2 = 4k^2$ for some integer k	basic algebra from 3, 4 (modus ponens)
formal proof – write as an argument and explicitly state each step $\boxed{\frac{n \text{ is even}}{\therefore n^2 \text{ is even}}}$	<ul> <li>6. if n<sup>2</sup> = 4k<sup>2</sup> for some integer k, then n<sup>2</sup> = 2(2k<sup>2</sup>) for that k</li> <li>7. n<sup>2</sup> = 2(2k<sup>2</sup>) for some integer k</li> <li>8. if n<sup>2</sup> = 2(2k<sup>2</sup>) for some integer k, then n<sup>2</sup> = 2k' for some integer k'</li> <li>9. n<sup>2</sup> = 2k' for some integer k', then n<sup>2</sup> is even</li> <li>11. n<sup>2</sup> is even</li> </ul>	basic algebra from 5, 6 (modus ponens) basic fact about integers from 7, 8 (modus ponens) definition of even
Proof. less formal – leave out explicit implications and instances of →	11. $n^*$ is even         Let n be an arbitrary integer.         1. n is even         2. $n = 2k$ for some integer k         3. $n^2 = 4k^2$ for that integer k         4. $n^2 = 2(2k^2)$ for that k         5. $n^2 = 2k'$ for some integer k'	trom 9, 10 (modus ponens) premise definition of even basic algebra basic algebra substituting $k' = 2k^2$
moaus ponens	<ol> <li>n<sup>2</sup> is even</li> <li>nce n was an arbitrary integer, the staten</li> </ol>	definition of even nent is true for all integers. $\Box$

# **Discovering Proofs – Advice**

- use the definition
  - e.g. a proof about even numbers likely needs to utilize the definition of "even"
- use the hypotheses
  - the hypotheses are the assumptions on which the conclusion will be based
  - previously-proved results can be used

#### Proof for all integers *n*, if *n* is even, then $n^2$ is even *Proof.* Let n be an arbitrary integer. 1. n is even premise less formal definition of even 2. n = 2k for some integer k leave out explicit 3. $n^2 = 4k^2$ for that integer k basic algebra implications and 4. $n^2 = 2(2k^2)$ for that k basic algebra instances of 5. $n^2 = 2k'$ for some integer k'substituting $k' = 2k^2$ modus ponens 6. $n^2$ is even definition of even Since n was an arbitrary integer, the statement is true for all integers. $\Box$ typical "formal" *Proof.* Let n be an arbitrary integer and assume n is even. Then n = 2kproof – written for some integer k by the definition of even, and $n^2 = 4k^2 = 2(2k^2)$ . Since in prose, may the product of integers is an integer, we have $n^2 = 2k'$ for some integer k'. not explicitly Therefore $n^2$ is even. Since n was an arbitrary integer, the statement is state all rules if true for all integers. reader can be expected to fill in those details CPSC 229: Foundations of Computation • Spring 2024

## **Discovering Proofs – Patterns**

- $\forall x P(x)$ 
  - tactic: let *x* be an arbitrary entity in the domain of discourse; show P(x) is true
  - showing P(x) for any entity means that it must hold for all entities
    don't use any facts about x beyond what is assumed
- ∃x P(x) existence proof
  - tactic: find an example a specific entity a for which P(a) is true
- disproving things
  - disprove  $\forall x P(x)$  with an example showing  $\exists x \neg P(x) a$  counterexample
  - disprove  $\exists x P(x)$  by showing  $\forall x \neg P(x)$

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### **Discovering Proofs – Patterns**

#### • $p \rightarrow q$ aka $\forall x ( P(x) \rightarrow Q(x) )$

- tactic: assume p, show q
  - $p \rightarrow q$  is true when p is false, so only need to handle the case where p is true
- tactic: show a chain of valid implications  $p \rightarrow r \rightarrow s \rightarrow \dots \rightarrow q$
- tactic: show the contrapositive  $(\neg q \rightarrow \neg p) indirect proof$
- p ^ q
  - tactic: show *p* and *q* separately
- p v q
  - tactic: assume  $\neg p$  and show q (based on  $p \lor q \equiv \neg p \to q$ )

#### • *p* ↔ *q*

- tactic: show  $p \rightarrow q$  and  $q \rightarrow p$
- in English: p if and only if q, p is necessary and sufficient for q
- tactic: show a chain of valid biconditionals  $p \leftrightarrow r \leftrightarrow s \leftrightarrow ... \leftrightarrow q$

4. Determine whether each of the following statements is true. If it true, prove it. If it is false, give a counterexample. a) Every prime number is odd. **b**) Every prime number greater than 2 is odd. c) If x and y are integers with x < y, then there is an integer z such that x < z < y. d) If x and y are real numbers with x < y, then there is a real number z such that x < z < y. 5. Suppose that r, s, and t are integers, such that r evenly divides s and s evenly divides t. Prove that r evenly divides t. **6.** Prove that for all integers n, if n is odd then  $n^2$  is odd. 7. Prove that an integer n is divisible by 3 iff  $n^2$  is divisible by 3. (Hint: give an indirect proof of "if  $n^2$  is divisible by 3 then n is divisible by 3.") 8. Prove or disprove each of the following statements. a) The product of two even integers is even. b) The product of two integers is even only if both integers are even. c) The product of two rational numbers is rational. d) The product of two irrational numbers is irrational. e) For all integers n, if n is divisible by 4 then  $n^2$  is divisible by 4. **f**) For all integers n, if  $n^2$  is divisible by 4 then n is divisible by 4. CPSC 229: Foundations of Computation • Spring 2024