Non-Regular Languages

Theorem 3.6. If L is a regular language, then there is some number n > 0such that any string w in L whose length is greater than or equal to n can be broken down into three pieces x, y, and z, w = xyz, such that

(i) x and y together contain no more than n symbols;

(ii) y contains at least one symbol;

(iii) xz is accepted by M

(xyz is accepted by M) xyyz is accepted by M etc.

show that { $a^nb^n \mid n \ge 0$ } is not regular

- let N be the threshold length and pick $a^{N}b^{N}$ as a string whose length is at least N
- show that a^Nb^N can't be written in the form xyz by showing that any choice for y that satisfies (i) and (ii) doesn't satisfy (iii)
 - since *xy* can't contain more than *N* symbols, both *x* and *y* contain only as
 - let k be the number of as in y since y can't be empty, $1 \le k \le N$
 - then $xz = a^{N+k}b^N$ which is not of the form a^nb^n and thus isn't accepted by by M

The Big Picture

Why do we care about being able to write computer programs that can recognize or generate languages?

- pattern matching
- L-systems
 - a system for describing fractal shapes
- compilers
 - being able to parse a program file
- ...

that DFAs can recognize the languages generated by regular expressions is good news for programs, but there are also languages, like { $a^nb^n | n \ge 0$ }, which aren't regular but are still easily recognizable by programs...





etc.

CPSC 229: Foundations of Computation • Spring 2024