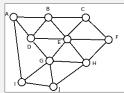
BFS – Implementation and Running Time bfs(G,s) for each vertex u in V-{s} do O(1) per element $\rightarrow O(n)$ total state[u] = "undiscovered" prev[u] = null state[s] = "discovered" enqueue, dequeue, isEmpty are O(1) prev[s] = null set, access state - want O(1), can do Q.enqueue(s) that with Map (hashtable) while Q is not empty do ← n repetitions u = Q.dequeue() } O(1) O(1) per element = for each edge (u,v) in 6.incidentEdges(u) do O(deg(u)) (adj list) if state[v] = "undiscovered" then O(n) (adj matrix) total over all iterations of state[v] = "discovered" while loop prev[v] = u • O(m) for adj list because the Q.enqueue(v) sum of the degrees is 2m state[u] = "processed" each edge is counted twice total over all iterations O(n²) for adj matrix because of while loop - O(n) while loop repeats n times O(n+m) for adjacency list implementation O(n2) for adjacency matrix implementation CPSC 327: Data Structures and Algorithms • Spring 2024

BFS-Based Algorithms

- unweighted shortest path
 - initialize dist[s] = 0, dist[v] = ∞ for all other vertices v
 - set dist[v] = dist[u]+1 when vertex v is discovered from vertex u
 - at the end, dist[v] has the length of the shortest path from s to v for all vertices v in the graph



intuition – all dist 1 vertices are discovered before dist 2, etc

BFS finds the unweighted shortest path because of how it traverses the graph

```
unweighted-shortest-path(G,s)
 for each vertex u in V-{s} do
   state[u] = "undiscovered"
   prev[u] = null
   dist[v] = ∞
 state[s] = "discovered"
 prev[s] = null
 dist[s] = 0
 Q.enqueue(s)
 while Q is not empty do
   u = Q.dequeue()
   for each edge (u,v) in G.incidentEdges(u) do
     if state[v] = "undiscovered" then
       dist[v] = dist[u]+1
       state[v] = "discovered"
       prev[v] = u
       Q.enqueue(v)
   state[u] = "processed"
```

BFS-Based Algorithms

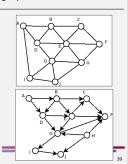
reachability

run bfs(s) – the vertices *reachable* from s are those marked as "processed" by bfs(s)

works with both undirected and directed graphs

intuition – we follow every edge leaving each discovered vertex, and every vertex put in the queue is eventually removed and marked as processed

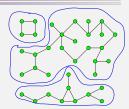
```
state[s] = "discovered"
Q.enqueue(s)
while Q is not empty do
    u = Q.dequeue()
for each edge (u,v) in G.incidentEdges(u) do
    if state[v] = "undiscovered" then
        state[v] = "discovered"
Q.enqueue(v)
state[u] = "processed"
```



BFS-Based Algorithms

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a graph with three connected components (circled)



connected components

 a connected component is a subgraph where every pair of vertices are connected by a path and there are no connections to other vertices not in the subgraph

```
c = 0
for each vertex v of graph G
if v has not been discovered
run bfs(v), setting comp[u] = c when each vertex u is
processed
C++
intuition - BFS finds all vertices
reachable from v along a path
```

BFS-Based Algorithms



bipartite graph detection / two-coloring

- on / bipartite not bipartite
- a bipartite graph is one whose vertices can be divided into two sets such that every edge connects a vertex in one set with a vertex in the other
- coloring refers to assigning labels (colors) to vertices so that no two adjacent vertices have the same label (color)
 - a two-coloring uses two colors

```
color[s] = 0
run bfs(s), setting color[v] = the opposite color of
color[u] for each discovery edge (u,v) and checking that
color[v] is the opposite color of color[u] for each non-
discovery edge (u,v)
```

if there is an edge (u,v) for which color[u] = color[v], the graph is not bipartite / two-colorable

intuition – following a path along discovery edges must alternate colors, since those edges are graph edges

- can't change the color of any vertex without changing them all
- _ non-discovery edges are also graph edges, and ends must be opposite colors 12

Takeaways

- · BFS algorithm
- BFS-based algorithms
 - graph traversal
 - reachability
 - unweighted shortest path
 - connected components
 - 2-coloring / detecting bipartite graphs

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4