

General Grammars

Definition 4.6. A *grammar* is a 4-tuple (V, Σ, P, S) , where:

1. V is a finite set of symbols. The elements of V are the non-terminal symbols of the grammar.
 2. Σ is a finite set of symbols such that $V \cap \Sigma = \emptyset$. The elements of Σ are the terminal symbols of the grammar.
 3. P is a set of production rules. Each rule is of the form $u \rightarrow x$ where u and x are strings in $(V \cup \Sigma)^*$ and u contains at least one symbol from V .
 4. $S \in V$. S is the start symbol of the grammar.
- a context-free grammar is a grammar where the rules are limited to the form $A \rightarrow x$ (a single non-terminal on the left)

General Grammars

- (general) grammars are more powerful than context-free grammars
 - we've seen several examples
- are there languages that can't be produced by grammars?
 - yes – for an alphabet Σ , there uncountably many languages but only countably many that can be generated by grammars

1. Find a derivation for the string $caabcb$, according to the first example grammar in this section. Find a derivation for the string $aabbcc$, according to the second example grammar in this section. Find a derivation for the string $aaaa$, according to the third example grammar in this section.

$S \rightarrow SABC$	$S \rightarrow SABC$	$S \rightarrow DTE$
$S \rightarrow \varepsilon$	$S \rightarrow X$	$T \rightarrow BTA$
$AB \rightarrow BA$	$BA \rightarrow AB$	$T \rightarrow \varepsilon$
$BA \rightarrow AB$	$CA \rightarrow AC$	$BA \rightarrow AaB$
$AC \rightarrow CA$	$CB \rightarrow BC$	$Ba \rightarrow aB$
$CA \rightarrow AC$	$XA \rightarrow aX$	$BE \rightarrow E$
$BC \rightarrow CB$	$X \rightarrow Y$	$DA \rightarrow D$
$CB \rightarrow BC$	$YB \rightarrow bY$	$Da \rightarrow aD$
$A \rightarrow a$	$Y \rightarrow Z$	$DE \rightarrow \varepsilon$
$B \rightarrow b$	$ZC \rightarrow cZ$	
$C \rightarrow c$	$Z \rightarrow \varepsilon$	

$$n_a(w) = n_b(w) = n_c(w)$$

$$a^i b^j c^i$$

$$a^{n^2}$$