Splay Trees Takeaways

- another form of restructuring operation
- randomized or heuristic approaches can result in good performance in practice because worst case scenarios are rare
- the notion of amortized analysis
 - an average based on the performance over a series of operations

Comparison

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- frequently-accessed elements are near the root in splay trees, resulting in faster access
 - advantageous in applications where there is *locality of reference* repeated access of related storage locations
- AVL trees are more tightly balanced than red-black trees, so faster retrieval but slower insertion and removal
 - "tightly balanced" \rightarrow smaller height
 - AVL trees are good for applications where trees are built once but searched often

Comparison

AVL trees and splay trees achieve O(log n) height by keeping the subtrees from getting too uneven.

2-4 trees and red-black trees achieve O(log n) height by constraining the maximum depth of any element.

- AVL trees, 2-4 trees, and red-black trees all have worstcase O(log n) operations
- splay trees have amortized O(log n) operations, with worst-case O(n) behavior

Comparison

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- splay trees have the simplest implementation

 just BST + restructuring operation no additional information to store/maintain
- red-black trees are more commonly used than 2-4 trees
 - easily built on top of binary trees just need to store a color bit
 - simpler to implement than 2-4 trees
- 'find' may restructure a splay tree
 - from a design perspective, having read-only operations change structure is undesirable

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 a consequence is that splay trees are not thread-safe for concurrent finds without extra bookkeeping

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Designing Data Structures

Studying balanced search trees reveals two tactics:

- it can be effective to add additional properties to the organization of the elements stored in order to improve runtime of an operation
 - e.g. AVL trees, 2-4 trees, red-black trees
- it can be effective (though harder to analyze) to piggyback local optimizations on other operations

 e.g. splay trees

In both cases, it is essential that the additional work does not overwhelm the savings gained.

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