

## HW 11

- give the steps of the dynamic programming process from class, not just a statement of the algorithm or a narrative account of your reasoning process
- establish the problem
  - specifications
  - examples
- identify avenues of attack
  - targets
  - paradigms and patterns
  - the series of choices
- define the algorithm
  - size
  - generalize / define subproblems: partial solution, alternatives, subproblem, memoization
  - base case(s)
  - main case
  - top level: initial subproblem, setup, wrapup
  - special cases
  - algorithm
- show termination and correctness
  - termination
  - correctness: establish the base case, show the main case, final answer
- determine efficiency
  - implementation
  - time and space
  - room for improvement
- only #2 was graded (or a different problem if you didn't hand in #2)
  - review those comments and consider their application to the other problems

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- #2 – the subproblem input in this case depends on the series of choices
  - process input – line break now or not?
    - the subproblem also requires the length of current line or the space left on current line – necessary for computing the cost (and for knowing if not breaking now is an option)
    - base case cost  $\geq 0$  because it is necessary to consider the cost of spaces on the current line if there are words on it
  - produce output – where is the next line break?
    - the subproblem only needs the remaining words because this decision is about the entire contents of the line – the cost of the space left can be computed from the length of the words between the current point and the next line break
    - base case cost is 0 because there are no words left to go on a line (the current line is empty)

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- avoid unnecessary work – even though the result may still be polynomial, it could be a better polynomial...
  - e.g. a `computeLineLength(i, j)` function to compute the cost of the remaining spaces for a line containing words  $i$  through  $j$  is a nice idea for organization, but consider its context

```
for j = i to n-1 do
  length ← computeLineLength(i, j)
  ...
```
  - this adds up the lengths of words  $i..i$ , then  $i..i+1$ ,  $i..i+2$ ,  $i..i+3$ , ... – a lot of repeated computation
  - instead incrementally add the length of word  $j$  to the running total in each iteration

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