

exec()

- runs a new program in the current process
 - loads the new program, overwriting the current code segment with the new program's code
 - re-initializes the heap and stack

exec()

```
int execl(const char *pathname, const char *arg, ...
          /* (char *) NULL */);
int execlp(const char *file, const char *arg, ...
          /* (char *) NULL */);
int execlx(const char *pathname, const char *arg, ...
          /*, (char *) NULL, char *const envp[] */);
int execv(const char *pathname, char *const argv[]);
int execvp(const char *file, char *const argv[]);
int execvpe(const char *file, char *const argv[],
            char *const envp[]);
```

- first argument (pathname or file) is the name of the executable file to run
- execv() vs execl() – execv takes an array of the program's arguments (argv), execl takes a list of separate parameters (arg0, arg1, arg2, ...)
 - argv[0], arg0 are the filename of the executable
- execvp(), execlp(), execvpe() – p means they use the shell's PATH environment variable to locate the executable

exec()

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <string.h>
#include <sys/wait.h>

int main(int argc, char *argv[]) {
    printf("hello (pid:%d)\n", (int) getpid());
    int rc = fork();
    if (rc < 0) { // fork failed; exit
        fprintf(stderr, "fork failed\n");
        exit(1);
    } else if (rc == 0) { // child (new process)
        printf("child (pid:%d)\n", (int) getpid());
        char *myargs[3];
        myargs[0] = strdup("wc"); // program: "wc"
        myargs[1] = strdup("p3.c"); // arg: input file
        myargs[2] = NULL; // mark end of array
        execvp(myargs[0], myargs); // runs word count
        printf("this shouldn't print out");
    } else { // parent goes down this path
        int rc_wait = wait(NULL);
        printf("parent of %d (rc_wait:%d) (pid:%d)\n",
              rc, rc_wait, (int) getpid());
    }
    return 0;
}
```

```
prompt> ./p3
hello (pid:29383)
child (pid:29384)
29 107 1030 p3.c
parent of 29384 (rc_wait:29384) (pid:29383)
prompt>
```

Motivating the API

- the separation of fork() and exec() allows the environment to be set up in the new process before the desired program is run
- an application of this is supporting redirection and pipes in the shell

```
prompt> wc p3.c > newfile.txt
```

- > redirects stdout to a file
- < redirects a file to stdin
- | sends stdout from one command to stdin for the next
 - wc p3.c | tee newfile.txt
 - sends the output from wc to tee, which echoes it to the screen and to newfile.txt

Implementing Redirection

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <string.h>
#include <fcntl.h>
#include <sys/stat.h>
#include <sys/wait.h>

int main(int argc, char *argv[]) {
    int rc = fork();
    if (rc < 0) {
        // fork failed
        fprintf(stderr, "fork failed\n");
        exit(1);
    } else if (rc == 0) {
        // child: redirect standard output to a file
        close(STDOUT_FILENO);
        open("./p4.output", O_CREAT|O_WRONLY|O_TRUNC,
            S_IRWXU);
        // now exec "wc"...
        char *myargs[3];
        myargs[0] = strdup("wc"); // program: wc
        myargs[1] = strdup("p4.c"); // arg: file to count
        myargs[2] = NULL; // mark end of array
        execvp(myargs[0], myargs); // runs word count
    } else {
        // parent goes down this path (main)
        int rc_wait = wait(NULL);
    }
    return 0;
}
```

```
prompt> ./p4
prompt> cat p4.output
32      109      846 p4.c
prompt>
```

the first available file descriptor is assigned when a file is opened – by closing stdout and immediately opening a new file, that new file will get the same file descriptor that had been assigned to stdout...and so things that think they are writing to stdout's are actually now writing to ./p4.output instead

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kill(), signal()

- `kill()` sends a signal to another process
 - SIGINT – generated by ctrl-C (term)
 - SIGSTOP – pause a process to be resumed later (stop)
 - SIGCONT – resume a paused process (cont)
 - SIGKILL – terminate a process immediately (term)
 - man 7 signal for a complete list
- `signal()` catches the signal and handles it
 - the *default disposition* defines what happens if the receiving process doesn't catch the signal

Term	Default action is to terminate the process.
Ign	Default action is to ignore the signal.
Core	Default action is to terminate the process and dump core (see <code>core(5)</code>).
Stop	Default action is to stop the process.
Cont	Default action is to continue the process if it is currently stopped.

- SIGKILL, SIGSTOP cannot be caught or ignored

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kill(), signal()

- users can generally only control their own processes
- the superuser (root) can control all processes